# EECC 730 VLSI Design Project Report $4 \times 4$ Signed/Unsigned Array Multiplier with BIST 

Aniket Dilip Mhatre<br>and<br>Sajin George<br>Instructor - Dr. Amlan Ganguly<br>TA - Anirudh Easwar

Submission Date - February 28, 2012

Top Level Directory: /shared/eecc630/eecc630-axmsxg/project_workspace/final_core Extra Credit - Universal 1 bit Multiplier Mode design(with 12 transistors) \& VHDL code to simulate universal $4 \times 4$ multipler with BIST

Rochester Institute of Technology
Department of Computer Engineering

Table of Contents

## 1. Abstract

2. Design Methodology and Theory
2.1 Mirror Adder
2.2 Minimizing delay Transistor Sizing:
2.3 Multiplier Core
2.4 BILBO
3. Results/Data Analysis
3.1 Results
3.2 Area of Layout
3.3 Number of NMOS and PMOS Transistors
3.4 Analysis of Result
4. Conclusion
5. Appendix

## 1. Abstract

The $4 \times 4$ universal multiplier with BIST was designed various blocks from different labs. The multiplier used mirrors full adders at different levels. A mode select line was introduced to set the sign bit for the operands. Build In Self Test circuit used was used to used to test the circuit. The BIST operates in the LFSR mode to generate a pseudo random sequence that is given as the inputs to the multiplier and the output of the multiplier is tested by the BIST. The multiplier can be even loaded the operands in parallel or serial shift in more.

## 2. Design Methodology and Theory

### 2.1 Mirror Adder

The design of the mirror adder is more compact and only use 24 transistors, this has been accomplished based on the observation that Sum can be factorized to reuse Cout as follows:

$$
\begin{aligned}
& S=A B C+(A+B+C) C o u t \\
& \text { Cout=(A.B)+(A xor B)Cin }
\end{aligned}
$$

The design is completely symmetrical, ie. The pMOS and nMOS circuits are exactly identical. The Critical Path goes from $C$ to Cout. The unit nMOS in the circuit has a resistance of $R$ and capacitance $C$ and unit pMOS has a resistance $2 R$ and capacitance $C$. the Capacitance is proportional to the width and the resistance is inversely proportional to the width. The Width of the transistors are selected to have equal rise and fall-time. The variation of the resistance is due to change in mobility of Holes and Electrons.

### 2.2 Minimizing delay Transistor Sizing:

The following techniques were used to reduce the delay of the circuit.
Feed the carry-in signal(C) to the inner inputs so that the internal capacitance is already discharged.(already done in the circuit diagram given)

- Make all transistors in the sum logic whose gate signals are connected to the carry-in and carry logic minimum size ( 1 unit e.g $4 \lambda$ ). This minimizes the branching effort on the critical path.
- The mirror adder has greater delay to compute Sum than Cout


### 2.3 D-FF

A D flip-flop is capable of storing the state of the input. It is a one bit memory element. The truth table of a D-FF is shown in Fig1. The circuit given in the lab is used to implement the D-FF for this lab. This circuit uses less number of transistors as compared to the Master Slave D-FF CMOS design. This circuit is suitable for high seed clocking as this circuit gets rid of the race condition. Figure 2 shows the circuit diagram of the D-FF. The D-FF is properly sized to obtain equal drive strength. The design is loaded with capacitor of 100 fF . Based on the schematic the
layout is designed and is optimized for minimum area. Post-layout and pre-layout delays are recorded after simulation.

### 2.3 MUX

The mux is used for taking the output from the XNOR gate and the output is given to the inverter of the OR gate. The mux is basically used to select inputs-either for LFSR mode or for user inputs through scan_in

### 2.4 Multiplier Core

Multiplication is an operation in which you scale one number by the other. It is an operation which involves series of additions. Bit wise multiplication for binary numbers is carried out in a similar fashion, however repeated addition introduces delay. To reduce the propagation delay the multiplication operation is broken into smaller operations to produce partial products. Further addition of the partial products yields the final product. A simplified equation for the multiplication of unsigned numbers is shown below.

$$
P=\left(\sum_{j=0}^{M-1} y_{j} 2^{j}\right)\left(\sum_{i=0}^{N-1} x_{i} 2^{i}\right)=\sum_{i=0}^{N-1} \sum_{j=0}^{M-1} x_{i} y_{j} 2^{i+j}
$$

[Lab Reference Material]

Figure 1 gives an example of $5 \times 5$ bit multiplication achieved using partial product.


Figure 1. $5 \times 5$ bit multiplication using partial products. [Lab Reference Material]

In order to achieve multiplication of signed as well as unsigned number two's complement multiplication is performed. Two's compliment multiplication takes into consideration the negative weights of the numbers. The equation for a two's complement $\mathrm{N} \times \mathrm{M}$ is given as:

$$
\begin{aligned}
P & =\left(-y_{M-1} 2^{M-1}+\sum_{j=0}^{M-2} y_{j} 2^{j}\right)\left(-x_{n-1} 2^{N-1}+\sum_{i=0}^{N-2} x_{i} 2^{i}\right) \\
& =\sum_{i=0}^{N-2} \sum_{j=0}^{M-2} x_{i} y_{j} 2^{i+j}+x_{N-1} y_{M-1} 2^{M+N-2}-\left(\sum_{i=0}^{N-2} x_{i} y_{M-1} 2^{i+M-1}+\sum_{j=0}^{M-2} x_{N-1} y_{j} 2^{j+N-1}\right)
\end{aligned}
$$

## [Lab Reference Material]

The multiplier speed depends on the number of partial products to be added. The modified Baugh-Wooley multiplier algorithm is used to reduce the number of partial products. The simplified technique to achieve two's complement $5 \times 5$ bit multiplication is shown below.

|  |  |  |  |  | $\begin{array}{r} y_{5} \\ x_{5} \\ \hline \end{array}$ | $\begin{array}{r} y_{4} \\ x_{4} \\ \hline \end{array}$ | $\begin{array}{r} y_{3} \\ x_{3} \\ \hline \end{array}$ | $\begin{array}{r} y_{2} \\ x_{2} \\ \hline \end{array}$ | $\begin{array}{r} y_{1} \\ x_{1} \\ \hline \end{array}$ | $\begin{array}{r} y_{0} \\ x_{0} \\ \hline \end{array}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | 1 | $\overline{x_{5} y_{0}}$ | $x_{0} y_{4}$ | $x_{0} y_{3}$ | $\mathrm{x}_{0} y_{2}$ | $x_{0} y_{1}$ | $\mathrm{x}_{0} \mathrm{y}_{0}$ |
|  |  |  |  | $\overline{x_{5} y_{1}}$ | $\mathrm{x}_{1} y_{4}$ | $\mathrm{x}_{1} \mathrm{y}_{3}$ | $\mathrm{x}_{1} \mathrm{y}_{2}$ | $\mathrm{x}_{1} y_{1}$ | $\mathrm{x}_{1} \mathrm{y}_{0}$ |  |
|  |  |  | $\overline{x_{5} y_{2}}$ | $\mathrm{x}_{2} y_{4}$ | $\mathrm{x}_{2} y_{3}$ | $\mathrm{x}_{2} \mathrm{y}_{2}$ | $x_{2} y_{1}$ | $x_{2} y_{0}$ |  |  |
|  |  | $\overline{x_{5} y_{3}}$ | $x_{3} y_{4}$ | $\mathrm{x}_{3} y_{3}$ | $\mathrm{x}_{3} \mathrm{y}_{2}$ | $\mathrm{x}_{3} \mathrm{y}_{1}$ | $\mathrm{x}_{3} y_{0}$ |  |  |  |
|  | $\overline{x_{5} y_{4}}$ | $x_{4} y_{4}$ | $x_{4} y_{3}$ | $\mathrm{x}_{4} \mathrm{y}_{2}$ | $x_{4} y_{1}$ | $\mathrm{x}_{4} y_{0}$ |  |  |  |  |
| 1 | $\mathrm{x}_{5} \mathrm{y}_{5} \overline{\mathrm{x}_{4} \mathrm{y}_{5}}$ | $\overline{x_{3} y_{5}}$ | $\overline{x_{2} y_{5}}$ | $\overline{x_{1} y_{5}}$ | $\overline{x_{0} y_{5}}$ |  |  |  |  |  |
| $p_{11}$ | $\mathrm{p}_{10} \quad \mathrm{P}_{9}$ | $\mathrm{P}_{8}$ | $\mathrm{p}_{7}$ | $p_{6}$ | $P_{5}$ | $p_{4}$ | $\rho_{3}$ | $\mathrm{p}_{2}$ | $\mathrm{p}_{1}$ |  |

[Lab Reference Material]

In order to achieve multiplication a mechanism is required to invert the input to certain partial products. This is done using the mode circuit. The block diagram for the modified BaughWooley multiplier is figure 2 . The shaded blocks represent the mode controlled blocks.


Figure 2. Modified Baugh-Wooley multiplier block diagram. [Lab Reference Material]

The multiplier design implemented in the project uses components implemented in the previous labs. The multiplier design was implemented in blocks. The block diagram of the multiplier core is shown below.


Figure 3. Block diagram of Multiplier
The block used were and_mirroradder (half adder and the AND gate) , and_fulladder (full adder and AND gate), mode, mirror adder, mode_fulladder (mode and fulladder) and AND gates. Full adder and half adder were implemented using full mirror adder. In order to implement half adder the cin of the full adder is connected to ground.

## MODE DESIGN (1 bit Universal Multiplier) :

The mode block was designed for extra credits and it was implemented using a different approach. The circuit of the mode block is shown in figure 4.


Figure. 4 Circuit Diagram of Mode Block

The logic behind the implementation is given as follows : $A$ and $B$ are given as inputs to the NAND gate to produce the output as /(A.B). The output of the NAND is given to an inverter to obtain (A.B). Now a MUX is used to select the inputs i.e. (A.B) and /(A.B) based on the select which is mode.

The challenge in the design was to implement the block in not more than 12 transistors. NAND uses 4 transistors, the inverter uses 2 inverter and the transmission gate MUX uses 6 transistors. Thus the block was implemented using 12 transistors. The mode block functionality is shown in the table below.

| MODE | Output |
| :--- | :--- |
| 0 | A.B |
| 1 | /(A.B) |

Table. MODE functionality

### 2.5 BILBO

Build in Logic Block Observer, is used to test the functionality of the multiplier design core. When chips are manufacture the chips have to be tested automatically for checking for errors. Hence there is a need for a testing circuit on all chips. The BILBO is a multipurpose chip which can be used to test the chip, generate random sequence, and load bits parally and test the signature.After loading the the bilbo with the required test vector we observe for the Golden Signature to test if the circuit function is proper. The BILBO block can work in different modes the LFSR mode, Shift Register Mode, Parallel shift in/ shift out mode and Serial Shift in/ Shift out mode. The diagram below show the basic schematic of a four bit bilbo [source: my courses lab handout]

fig 2.2.1
the various modes of operation of the bilbo is controlled my the control line $\mathrm{C} 1, \mathrm{C} 2, \mathrm{C} 3$. table 2.2.1 show the various modes of operation.

| C1 | C2 | C3 | Mode |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | Shift - register |
| 0 | 0 | 1 | LFSR MODE |
| 0 | 1 | $X$ | Load " 0 " |
| 1 | 0 | 1 | Signature - Analysis |
| 1 | 1 | $X$ | Parallel Mode |
| 1 | 0 | 0 | Not used |

table 2.2.1
Modes of Operation[source lab-hand]

## Shift- Register Mode:

In the shift registered mode is used to load the values into the $D$ filpflop serially, this process of loading the bits will require 8 clock cycles. The data to be fed in is given at the SHIFTIN pin.

(a)

## LFSR mode:

LFSR ( Liner Feedback Shift Register) is used to create pseudo random sequence that will create a pseudo random input sequence to the testing core. The pseudo output of the LFSR can be seen in the (appendix) The initial value of the inside the LFSR is called as the seed value.


## Signature Analysis

The signature analysis is used to test the output from the the testing core. The output of the multiplier is fed into the BILBO us int the T1-T8, parallel input lines of the bilbo.
$C_{1}=1, C_{2}=0, C_{3}=q_{i}, 1$ Signature-Analysis Mode


## Parallel Mode

The Parallel mode is used to load values into the D flipflop in parallel.


## 3. Results/Data Analysis

### 3.1 Results

Multiplier:

|  | Vdd(V) | Temp(C) | Load | Rise Time | Fall Time | Tp,HL | Tp,LH |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Worst Case(Pre-layout) | 1.62 | 125 | 100 | 105.32 ps | 102.99 ps | 2.516 Ns | 2.138 Ns |
| Worst Case (post Layout) | 1.62 | 125 | 100 | 1.58 ns | 880.9 ps | 2.6629 | 2.0107 |

## Multiplier with BIST:

|  | Vdd(V) | Temp(C) | Load | Rise Time | Fall Time | Tp,HL | Tp,LH |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Worst Case(Pre-layout) | 1.62 | 125 | 100 | 971.87 | 223.91 | 39.979 ns | 39.924 |
| Worst Case (post Layout) | 1.62 | 125 | 100 | 5.25 ns | 1.14 ns | 39.977 |  |
|  |  |  |  |  |  |  | 39.841 |

## Functionality Check:

The functionality of the circuit was tested using the following test vectors and output verified.

| Test Vector | A | B | Product <br> (Unsigned) | Product <br> (Signed) |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 0000 | 0000 | 00000000 | 00000000 |
| 2 | 1111 | 1111 | 11100001 | 00000001 |
| 3 | 0101 | 0101 | 00011001 | 00011001 |
| 4 | 1010 | 1010 | 01100100 | 00100100 |
| 5 | 0000 | 1111 | 00000000 | 00000000 |
| 6 | 1111 | 0000 | 00000000 | 00000000 |
| 7 | 0000 | 0000 | 00000000 | 00000000 |

### 3.2 Area of Layout

Multiplier core area: $1470 \times 535 \lambda$
BIST area: $1465 \times 150 \lambda$
Entire area (Core + BIST) : $1470 \times 860 \lambda$
Fraction of pad area used: 8.54 \%

### 3.3 Number of NMOS and PMOS Transistors

Multiplier Core : 456 transistors
BILBO : 174 transistors

Entire core (BILBO+ Multiolier Core) : 630 transistors

### 3.4 Analysis of Result

From the result it is clearly visible that there is a one clock cycle delay in the output of the signature analyser because the load shift register shift for every clock, and since all the D flip flops are at the initial state of " 0 " it takes one more cycle to for outputs to change at the output.

The signature was verified by manually calculating the golden signature for the the test vectors, A VHDL program (appendix 2.3) was written to simulate the whole system and the output of the the VHDL code matched the output of the $4 \times 4$ multiplier with BIST.

The output of the post layout Bilbo was different due to the additional parasitic capacitance. This further increases the delay caused by the D Flip Flop.

## 4. Conclusion

The $4 \times 4$ universal multiplier with, BIST was designed using a hierarchical design. The POST layout performance was lower compared to the PRE- LAYOUT due to the play of parasitic capacitance. there is increase of $10 \%$ in the delay, which shows the efficiency of the design. The golden signature for the test vector was verified and a the obtained output was verified againts a VHDL simulation of the whole circuit.
5. Appendix

Layout:

and


Bilbo


Core


Dff


Mirror adder


MODE


MUX


Multiplier Core layout

Schematic:


And


Bilbo


Mirror adder



Multiplier Core

Waveforms:

## Multiplier Core wave forms

## Post-layout





## Pre-layout



## CORE:

Pre-layout


## Post-layout



LVS:

## \#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

```
## ##
```

\#\# CALIBRE SYSTEM \#\#
\#\# \#\#
\#\# LVS REPORT \#\#
\#\# \#\#
\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#\#

REPORT FILE NAME: /shared/eecc630/eecc630axmsxg/project_workspace/final_core_beta/lvs/lvs.rep LAYOUT NAME:
/shared/eecc630/eecc630-
axmsxg/project_workspace/final_core_beta/layout/final_core_beta
SOURCE NAME:
/shared/eecc630/eecc630-
axmsxg/project_workspace/final_core_beta/lvs
RULE FILE: /tools/mentor/adk/technology/ic/process/tsmc018.rules
LVS MODE: Mask
RULE FILE NAME: /tools/mentor/adk/technology/ic/process/tsmc018.rules
CREATION TIME: Tue Feb 28 13:46:31 2012
CURRENT DIRECTORY: /shared/eecc630/eecc630-axmsxg/project_workspace
USER NAME: axm9958


## INITIAL NUMBERS OF OBJECTS

## Layout Source Component Type



Ports: 3030

Nets: 594589 *

Instances: 696566 * mn (4 pins)
696566 * mp (4 pins)
$\qquad$
Total Inst: 13921132

NUMBERS OF OBJECTS AFTER TRANSFORMATION

## Layout Source Component Type

------ ------

Ports: 3030

Nets: 589589

Instances: 566566 mn (4 pins)
$566566 \mathrm{mp}(4$ pins $)$
$\qquad$
Total Inst: 11321132

* = Number of objects in layout different from number in source.
o LVS Setup:
LVS COMPONENT TYPE PROPERTY $\quad$ phy_comp element comp
LVS COMPONENT SUBTYPE PROPERTY model
LVS PIN NAME PROPERTY
LVS POWER NAME
phy_pin
LVS GROUND NAME


| SOURCE CASE NO | NO |
| :---: | :---: |
| LVS COMPARE CASE NO | NO |
| LVS DOWNCASE DEVICE | NO |
| LVS REPORT MAXIMUM | 50 |
| LVS PROPERTY RESOLUTION MAXIMUM 32 |  |
| // LVS SIGNATURE MAXIMUM |  |
| // LVS FILTER UNUSED OPTION |  |
| // LVS REPORT OPTION |  |
| LVS REPORT UNITS YES |  |
| // LVS NON USER NAME PORT |  |
| // LVS NON USER NAME NET |  |
| // LVS NON USER NAME INSTANCE |  |
| // Reduction |  |
| LVS REDUCE SERIES MOS | NO |
| LVS REDUCE PARALLEL MOS | YES |
| LVS REDUCE SEMI SERIES MOS | S NO |
| LVS REDUCE SPLIT GATES | YES |
| LVS REDUCE PARALLEL BIPOLAR | AR YES |
| LVS REDUCE SERIES CAPACITORS | ORS YES |
| LVS REDUCE PARALLEL CAPACITORS | CITORS YES |
| LVS REDUCE SERIES RESISTORS | RS YES |
| LVS REDUCE PARALLEL RESISTORS | TORS YES |
| LVS REDUCE PARALLEL DIODES | ES YES |
| LVS REDUCTION PRIORITY | PARALLEL |
| // Filter |  |

// Filter

LVS FILTER sch_filter_direct_open OPEN SOURCE DIRECT
LVS FILTER sch_filter_direct_short SHORT SOURCE DIRECT

LVS FILTER sch_filter_mask_open OPEN SOURCE MASK
LVS FILTER sch_filter_mask_short SHORT SOURCE MASK
LVS FILTER lay_filter_direct_open OPEN LAYOUT DIRECT
LVS FILTER lay_filter_direct_short SHORT LAYOUT DIRECT
LVS FILTER v OPEN
LVS FILTER i OPEN
LVS FILTER e OPEN
LVS FILTER f OPEN
LVS FILTER g OPEN
// Trace Property
// TRACE PROPERTY mn instpar(w) width w 0
// TRACE PROPERTY mp instpar(w) width w 0
// TRACE PROPERTY me instpar(w) width w 0
// TRACE PROPERTY md instpar(w) width w 0
// TRACE PROPERTY mn instpar(I) length IO
// TRACE PROPERTY mp instpar(I) length IO
// TRACE PROPERTY me instpar(I) length IO
// TRACE PROPERTY md instpar(I) length IO
// TRACE PROPERTY r instpar(r) resistance r 0
// TRACE PROPERTY c instpar(c) capacitance c 0
// TRACE PROPERTY d instpar(a) area a 0
// TRACE PROPERTY d instpar(p) perimeter p 0

o Statistics:

5 isolated layout nets were deleted.

390 layout mos transistors were reduced to 130.
260 mos transistors were deleted by parallel reduction.
o Isolated Layout Nets:
(Layout nets which are not connected to any instances or ports).

36(-253.000,-443.000) 37(917.000,261.000) 41(-284.000,268.000) 46(481.000,-282.000) 47(946.000,-436.000)
o Initial Correspondence Points:

Ports: VDD GND MODE SHIFTIN T2 T3 T4 T5 T6 T7 T8 CLK SHIFTOUT P2 P1 P0 P4 P6 P5 P3 SHIFTIN2 T1 C12 C11 C13 C22 C21 COUT RST C32


Total CPU Time: 0 sec
Total Elapsed Time: 0 sec

